CS355: Applied Zero Knowledge Proofs

Spring 2025

Assignment #1

Due: 11:59pm on Wed, Apr. 16, 2025, on Gradescope (each answer on a separate page)

Problem 1. (*ZK proof variants*) In Lecture 2 we saw a classic zero-knowledge interactive proof for the 3-coloring relation. For a graph G = (V, E), in every iteration, the verifier chose a uniform random edge e = (i, j) in *E*. We explained in the lecture that if the verifier instead chose a random pair $(i, j) \in V^2$ then the proof system would not be zero knowledge. Suppose that the verifier always chose its query as a random edge e = (i, j) in *E'*, where *E'* is a proper subset of *E* (i.e., $E' \subset E$ but $E' \neq E$). The rest of the protocol is unchanged. You can assume that the set *E'* is part of the statement, that is the prover and verifier are given (G, E').

- a. Is the resulting proof system HVZK? If so, show a simulator.
- **b.** Is the resulting (repeated) proof system sound? If so explain why. If not, show a statement (G, E') where G is not 3-colorable, but the verifier always accepts.

Problem 2. (Sigma protocols) Let G be a cyclic group of prime order q with generator $g \in G$. Consider the following relation

$$R := \left\{ \left((u, v, w \in G), (\alpha \in \mathbb{Z}_q) \right) : u = g^{\alpha}, v = g^{(\alpha^2)}, w = g^{(\alpha^3)} \right\}$$

- **a.** Use the generalized Schnorr protocol to write out an explicit HVZK proof system for R. Hint: consider the homomorphism $\varphi(\gamma) := (g^{\gamma}, u^{\gamma}, v^{\gamma})$ from \mathbb{Z}_q to G^3 .
- **b.** What is the transcript of this protocol?
- c. Show that your proof system from part (a) is complete.
- **d.** Give an explicit proof that your proof system from part (a) is HVZK. That is, construct a simulator S(u, v, w) that simulates the transcript whenever $(u, v, w) \in L(R)$.
- e. Prove that your proof system from part (a) is knowledge sound. It suffices to show that your proof system from part (a) is 2-special sound.